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MATHEMATICS[www.elsevier.com/locate/disc](http://www.elsevier.com/locate/disc)On minimal cutsets in  $P_5$ -free minimal imperfect graphsVincent Barré<sup>a,\*</sup>, Jean-Luc Fouquet<sup>b</sup><sup>a</sup>*Département Services et Réseaux de Communication, IUT de Laval, Université du Maine,  
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**Abstract**

Let  $G$  be a minimal imperfect  $P_5$ -free graph (i.e. a minimal imperfect graph not containing a path on 5 vertices as induced subgraph) and let  $S$  be a minimal cutset of  $G$ . In this paper we study several properties of such cutsets, in particular we prove that the subgraph induced by  $S$  is connected, contains a  $P_4$ , cannot induce a bipartite subgraph of  $G, \dots$ . As a by-product we also give a structural characterization of such graphs. © 2001 Elsevier Science B.V. All rights reserved.

**1. Introduction**

A graph is *perfect* if the vertices of any induced subgraph  $H$  can be colored, in such a way that no two adjacent vertices receive the same color, with a number of colors (denoted by  $\chi(H)$ ) not exceeding the cardinality  $\omega(H)$  of a maximum clique of  $H$ . For a graph  $G$  we denote by  $\alpha(G)$  the cardinality of a maximum stable set of  $V(G)$ .

A graph is *minimal imperfect* if all its proper induced subgraphs are perfect but it is not. In particular,  $\omega(G) + 1 = \chi(G)$  for  $G$  minimal imperfect. All the notions not defined here may be found in [2].

It is an easy task to check that an odd chordless cycle of length at least five (usually called a *hole*), as well as its complement (usually called an *anti-hole*) are minimal imperfect graphs.

The remark above and some early results concerning perfect graphs determined Berge [1] to formulate the two following conjectures (known as the Strong and the Weak Perfect Graph Conjecture)

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- (SPGC) A graph is perfect if and only if it does not contain an odd hole or an odd anti-hole as an induced subgraph.
- (WPGC) A graph is perfect if and only if its complement is perfect.

While the SPGC is still unsettled, the WPGC is an easy consequence of the following theorem of Lovász [8]:

**Theorem 1** (The Perfect Graph Theorem). *A graph  $G = (V, E)$  is perfect if and only if for every induced subgraph  $H$  of  $G$  the following inequality holds:*

$$\alpha(H) \cdot \omega(H) \geq |H|.$$

This theorem was the first step towards a new approach of minimal imperfect graphs. We can deduce from this theorem, that in a minimal imperfect graph  $G$ ,

- (1)  $n = \alpha \cdot \omega + 1$ .
- (2) For every vertex  $v \in V(G)$ ,  $G - v$  has a unique partition into  $\alpha\omega$ -cliques (i.e. a clique of size  $\omega$ ) and a unique partition into  $\omega$   $\alpha$ -stable sets.
- (3) Each vertex of  $G$  is in exactly  $\alpha$   $\alpha$ -stable sets and in exactly  $\omega$   $\omega$ -cliques.

Bland et al. [3] defined a graph to be *partitionable* if there exist two numbers  $\alpha, \omega \geq 2$  such that (1) and (2) hold. Further refinements along this line are due to Padberg [13]. We only need the following property:

**Proposition 2.** *In a minimal imperfect graph  $G$ , given two vertices  $u$  and  $v$ , there exists an  $\omega$ -clique containing  $u$  and not containing  $v$ .*

**Proof.** Use property (3).  $\square$

Two of the most useful graphical properties of minimal imperfect graphs were found by Meyniel and Chvátal. To describe their results we need introduce a few definitions.

Two nonadjacent vertices  $x, y$  form an *even pair* if all chordless paths joining  $x$  to  $y$  have an even number of edges. A set  $C$  of vertices of a connected graph  $G$  is called a *star-cutset* if  $G - C$  is not connected and there is a vertex  $x$  in  $C$  adjacent to all other vertices of  $C$ ;  $x$  is called the *center* of  $C$ .

**Lemma 3** (Meyniel [12]). *No minimal imperfect graph contains an even pair.*

**Lemma 4** (Chvátal [4]). *No minimal imperfect graph contains a star-cutset.*

Now, we denote by  $N_G(x)$  the set of vertices of  $G$  adjacent to  $x$ ; when there can be no confusion we shall write  $N(x) = N_G(x)$ .

**Lemma 5.** *No minimal imperfect graph contains two nonadjacent vertices  $u$  and  $v$  such that*

$$N(v) \subseteq N(u) \quad \text{or} \quad N(u) \subseteq N(v)$$

(one says they have comparable neighbourhoods).

**Proof.** Assume, for example, that  $N(v) \subseteq N(u)$  then  $S = \{u\} \cup N(v)$  is a star-cutset except if  $V = \{u, v\} \cup N(u)$  and  $N(u) = N(v)$  (because  $|V \setminus S| = 1$ ) but in this case  $\{u, v\}$  form an even-pair, which is impossible.  $\square$

We now introduce another useful tool.

**Definition 6.** Let  $G = (V, E)$  be a minimal imperfect graph and let  $u, v$  be two nonadjacent vertices of  $G$ . We denote by  $G + uv$  the graph  $(V, E \cup \{uv\})$  and one says that  $u, v$  is a *co-critical* pair if:  $\omega(G + uv) = \omega(G) + 1$ . Analogously, an edge  $xy \in E$  is said to be *critical* if removing this edge from  $E$  increases the stability number of  $G$ .

Giles et al. [7] (see also Markossian and Karapetian [11]) prove this interesting result:

**Theorem 7** (Giles et al. [7], Markossian and Karapetian [11]). *No minimal imperfect graph (different from an odd hole or an odd anti-hole) contains a cycle of critical edges.*

We are interested in the following consequence of their result :

**Lemma 8.** *No minimal imperfect graph (different from an odd hole or an odd anti-hole) contains three vertices inducing a stable set of co-critical pairs.*

There are many conjectures concerning minimal imperfect graphs, in particular Chvátal [5] proposes:

**Conjecture 9.** Every minimal imperfect graph  $G$  containing no odd hole and no odd antihole has the following properties:

- (1) For each cutset  $C$  of  $G$ , the subgraph  $[C]_G$  of  $G$  induced by  $C$  is connected.
- (2) For each cutset  $C$  of  $G$ , the subgraph  $[C]_G$  contains a  $P_4$ .

The main purpose of this paper is to prove this conjecture for  $P_5$ -free graphs. As a by-product we obtain a structural characterization of minimal imperfect and  $P_5$ -free graphs.

## 2. On minimal cutsets

We shall study properties of minimal cutsets in minimal imperfect  $P_5$ -free graphs. We call the *complete join* of two (disjoint) graphs  $A = (V(A), E(A))$  and  $B = (V(B), E(B))$  the graph with the vertex set  $V(A) \cup V(B)$  and the edge set  $E(A) \cup E(B) \cup \{ab | a \in V(A), b \in V(B)\}$ . A graph is called *Berge* if it contains no odd hole and no odd antihole. We will prove the following results in part 4 of the paper:

**Lemma 10.** *Let  $G$  be a minimal imperfect and  $(P_5, C_5)$ -free graph, and let  $S$  be a minimal cutset of  $G$ . If  $[S]_G$  is the complete join of two graphs  $A$  and  $B$  (with  $V(A) \neq \emptyset$  but  $V(B)$  may be empty) then  $A$  is isomorphic to a connected subgraph of  $G$ .*

**Remark 11.** When we perform the complete join of a graph  $A$  and a graph  $B$  such that  $V(A) \neq \emptyset$  and  $V(B) = \emptyset$ , the result is the graph  $A$ . So, a minimal cutset  $S$  of  $G$  induces a connected subgraph of  $G$ , else we can choose  $V(A) = S$  and  $V(B) = \emptyset$  which contradicts Lemma 10.

So, we have

**Theorem 12.** *If  $G$  is a minimal imperfect and  $(P_5, C_5)$ -free graph then for each minimal cutset  $S$  of  $G$ , the subgraph of  $G$  induced by  $S$  is connected.*

**Theorem 13.** *If  $G$  is a minimal imperfect and  $(P_5, C_5)$ -free graph then for each minimal cutset  $S$  of  $G$ , the subgraph of  $G$  induced by  $S$  contains a  $P_4$ .*

We know that, in a minimal imperfect graph, the neighbourhood of any vertex is a minimal cutset. So, we have the following corollaries:

**Corollary 14.** *If  $G$  is a minimal imperfect and  $(P_5, C_5)$ -free graph, then for each vertex  $v$ , the subgraph induced by  $N(v)$  is connected.*

**Remark 15.** One can show that this property is also true if the graph  $G$  is partitionable and  $(P_5, C_5)$ -free.

**Corollary 16.** *If  $G$  is a minimal imperfect and  $(P_5, C_5)$ -free graph, it cannot have a vertex  $v$  such that the subgraph of  $G$  induced by  $N(v)$  is  $P_4$ -free.*

**Corollary 17.** *If  $G$  is a  $(P_5, C_5)$ -free graph such that every its induced subgraph  $H$  has a vertex  $v$  such that the subgraph of  $H$  induced by  $N_H(v)$  is  $P_4$ -free then  $G$  is perfect.*

**Remark 18.** This class of perfect graphs is a subclass of Slightly Triangulated Graphs the perfection of which was shown in [10].

We can also derive from Lemma 10 the following characterization of minimal imperfect and  $(P_5, C_5)$ -free graphs. Let  $G = (V, E)$  be such a graph, let  $v$  be a vertex of  $G$  and let  $M(v) = M_G(v) = N_G(v) = \{u \in V \mid u \neq v, uv \notin E\}$ . For  $w \in N(v)$  we put  $\mathcal{M}(w) = M(v) \cap N_G(w)$ . For a subset  $Y$  of  $V$  we write  $N_Y(u) = N(u) \cap Y$ .

**Theorem 19.** *For every vertex  $v \in V(G)$ , one has a partition of  $V(G)$  in*

$$\{v\}, W, Y, \text{ and } M$$

*such that  $N(v)$  is partitioned in  $W$  and  $[Y]_G$  (with  $Y \neq \emptyset$ ,  $[Y]_G$  is connected and  $W = \{w\} \cup N_{N(v)}(w)$ , where  $w$  is a vertex such that  $M = \mathcal{M}(w)$ ).*

Now, let  $\mathcal{B}$  be the family of bipartite graphs and let  $\mathcal{B}^*$  be the family defined by

- $\mathcal{B} \subseteq \mathcal{B}^*$ .
- $\forall G_1, G_2 \in \mathcal{B}^*$ , the complete join and the disjoint union of  $G_1$  and  $G_2$  are in  $\mathcal{B}^*$ .

Gallai [6] proved that, for each vertex  $v$  of a minimal imperfect graph  $G$ , the neighbourhood of  $v$  induces a connected subgraph of the complement of  $G$ ; hence if  $[N(v)]_G \in \mathcal{B}^*$ , then we have  $[N(v)]_G$  is disconnected or  $N(v)$  induces a bipartite graph (the latter case is impossible if  $G$  is Berge because this implies that  $\omega(G) = 3$  and Tucker [15] has shown that SPGC is true for  $K_4$ -free graphs).

Then, if Lubiw's conjecture [9] (in  $G$  minimal imperfect and Berge,  $\forall v \in V(G)$ ,  $[N(v)]_G$  is connected) is true, for each vertex  $v$  of  $G$  we have  $[N(v)]_G \notin \mathcal{B}^*$ . So, we propose this weaker conjecture:

**Conjecture 20.** *If  $G$  is a minimal imperfect Berge graph then for every vertex  $v$  of  $G$  we have  $[N(v)]_G \notin \mathcal{B}^*$ .*

This conjecture implies that  $[N(v)]_G$  contains a  $P_4$ . In the case of minimal imperfect and  $P_5$ -free graphs, we prove this conjecture for every minimal cutset. For doing this, our key lemma is the following:

**Lemma 21.** *Let  $G$  be a  $P_5$ -free minimal imperfect Berge graph and let  $S$  be a minimal cutset of  $G$ . If  $[S]_G$  is the complete join of two graphs  $A$  and  $B$  ( $V(A) \neq \emptyset$ , but  $V(B)$  may be empty), then  $A$  is not isomorphic to a bipartite subgraph of  $[S]_G$ .*

### 3. Preliminary results

Let  $G = (V, E)$  be a minimal imperfect and  $(P_5, C_5)$ -free graph. Let  $v \in V$  be a vertex of  $G$ .

**Lemma 22.** *Let  $G$  be a  $(P_5, C_5)$ -free graph. Let  $v \in V$  and  $x, y \in N(v)$  such that  $xy \notin E$  then  $\mathcal{M}(x) \subseteq \mathcal{M}(y)$  or  $\mathcal{M}(y) \subseteq \mathcal{M}(x)$  (if  $\mathcal{M}(x)$  and  $\mathcal{M}(y)$  are nonempty).*

**Proof.** Suppose that there exist  $x, y \in N(v)$  contradicting the hypothesis. So, there exist  $a \in \mathcal{M}(x) \setminus \mathcal{M}(y)$  and  $b \in \mathcal{M}(y) \setminus \mathcal{M}(x)$ , but  $\{a, x, v, y, b\}$  induces a  $P_5$  or a  $C_5$ , a contradiction.  $\square$

Now, let  $S$  be a minimal cutset of  $G$ , and let  $G_1, G_2, \dots, G_q$  ( $q \geq 2$ ) be the connected components (meant as maximal subsets of vertices inducing connected subgraphs) of  $G \setminus S$ . We say that a vertex  $x \in V$  is *incomplete* for a set  $A$  if we can find a vertex  $y \in A$  such that  $xy \notin E$  otherwise vertex  $x$  is said to be *complete* for  $A$ .

**Lemma 23.** *For every vertex  $v \in S$ ,  $v$  is incomplete for, at most, one component  $G_i$  ( $1 \leq i \leq q$ ).*

**Proof.** Suppose that  $v \in S$  is incomplete for two components  $G_i$  and  $G_j$  ( $i \neq j$ ). Each component is connected and then, one can find a vertex  $v_i$  in  $N_{G_i}(v) \neq \emptyset$  and a vertex  $w_i$  in  $V(G_i) \setminus N_{G_i}(v) \neq \emptyset$  such that  $v_i w_i \in E$  (we can do this for  $G_j$ ) but then the subset  $\{w_i, v_i, v, v_j, w_j\}$  induces a  $P_5$ .  $\square$

**Remark 24.** Let  $v$  be a vertex of  $G_i$  ( $1 \leq i \leq q$ ) then for every vertex  $x$  in  $S$  we have  $\mathcal{M}(x) \neq \emptyset$  ( $G_j \subseteq \mathcal{M}(x)$  for  $j \neq i$ ). So from now on, we will not verify if  $\mathcal{M}(x)$  and  $\mathcal{M}(y)$  are nonempty when applying Lemma 22.

#### 4. The proofs

**Proof of Lemma 10.** Let  $G$  be a minimal imperfect and  $(P_5, C_5)$ -free graph, let  $S$  be a minimal cutset of  $G$  and let  $G_1, G_2, \dots, G_q$  ( $q \geq 2$ ) be the connected components of  $G \setminus S$ . Suppose that  $[S]_G$  is the complete join of two graphs  $A$  and  $B$  ( $V(A) \neq \emptyset$ , and  $V(B)$  may be empty) such that  $A$  is isomorphic to a disconnected subgraph (we denote it also by  $A$ ) of  $[S]_G$ . Let  $A_1, A_2, \dots, A_p$  ( $p \geq 2$ ) be the connected components of  $A$ .

*Case 1:*  $q \geq 3$ . Let  $x \in V(A)$  be a vertex such that  $N_G(x) - S$  is maximal. We can suppose, without loss of generality, that  $x \in A_1$ . Then for every vertex  $y$  in  $A_2$ , the neighbourhood of  $y$  in  $\bigcup_{1 \leq i \leq q} G_i$  is included in the neighbourhood of  $x$  in  $\bigcup_{1 \leq i \leq q} G_i$ . Indeed, since  $q \geq 3$  there exists one index  $k$  ( $1 \leq k \leq q$ ) such that  $x$  and  $y$  are complete for  $G_k$  (Lemma 23). Let  $v \in G_k$ , we have  $\bigcup_{i \neq k} G_i \subseteq M(v)$  so if we apply Lemma 22 with this vertex  $v$ , we have  $\mathcal{M}(x) \subseteq \mathcal{M}(y)$  or  $\mathcal{M}(y) \subseteq \mathcal{M}(x)$  and, particularly,

$$N_{\bigcup_{i \neq k} G_i}(x) \subseteq N_{\bigcup_{i \neq k} G_i}(y) \quad \text{or} \quad N_{\bigcup_{i \neq k} G_i}(y) \subseteq N_{\bigcup_{i \neq k} G_i}(x),$$

that is  $N_{\bigcup_i G_i}(x) \subseteq N_{\bigcup_i G_i}(y)$  or  $N_{\bigcup_i G_i}(y) \subseteq N_{\bigcup_i G_i}(x)$  since  $x$  and  $y$  are complete for  $G_k$ . But  $N_{\bigcup_i G_i}(x)$  is maximal, therefore  $N_{\bigcup_i G_i}(y) \subseteq N_{\bigcup_i G_i}(x)$ .

If  $x$  is incomplete for a component  $G_i$  then  $\{x\} \cup N_G(x)$  forms a star disconnecting subgraphs induced by subsets  $A_2$  and  $G_i \setminus N_{G_i}(x)$ . So, we can suppose that  $x$  is complete for  $G_1, G_2, \dots, G_q$  but in this case, if  $V(A) \neq \{x\} \cup A_2$ , the subset  $\{x\} \cup G_1 \cup G_2 \cup \dots \cup G_q \cup V(B)$  forms a star cutset. Then we only need to study the following case:

**Claim 25.** *If  $[S]_G$  is the complete join of two graphs  $A$  and  $B$  then we cannot have*

- $V(A) = \{x\} \cup Y$  (where  $A$  is disconnected and the subgraph of  $A$  induced by  $Y$  is connected), and
- $x$  complete for  $G_1 \cup \dots \cup G_q$  ( $q \geq 2$ ).

The subset  $Y$  can be partitioned into  $U, Y_1, \dots, Y_q$  where

- $U$ : vertices in  $Y$  which are complete for  $G_1, \dots, G_q$ ;
- $Y_i$  ( $1 \leq i \leq q$ ): vertices in  $Y$  which are incomplete for  $G_i$ .

We first show that  $U = \emptyset$ . We assume that  $y \in U \subseteq Y$  is complete for  $G_1, \dots, G_q$  ( $q \geq 2$ ) and show that this leads to a contradiction:

- If  $Y - \{y\} \neq \emptyset$  then  $\{y\} \cup G_1 \cup \dots \cup G_q \cup V(B)$  forms a star-cutset.
- If  $Y - \{y\} = \emptyset$  then  $x$  and  $y$  have the same neighbourhood (which contradicts Lemma 5).

Now, assume there exist two indices  $i$  and  $j$  ( $i \neq j$ ) such that  $Y_i \neq \emptyset$  and  $Y_j \neq \emptyset$ . Let  $u \in Y_i$ ,  $v \in Y_j$ ,  $t \in G_j \setminus N(v) \neq \emptyset$  and  $r \in G_i \setminus N(u) \neq \emptyset$  then the subset  $\{u, t, x, r, v\}$  induces a  $P_5$  or a  $C_5$ . So, we can suppose, without loss of generality, that  $Y_1 \neq \emptyset$  and that for each  $j$ ,  $2 \leq j \leq q$ ,  $Y_j = \emptyset$  (i.e.  $Y = Y_1$ ). Moreover, no vertex  $y$  in  $Y$  is *universal* for  $Y$  (that is adjacent to all vertices in  $Y$ ), otherwise

- If  $\omega([G_2]_G) = \omega - 1$ , then any  $\omega$ -clique containing  $v \in G_2$  also contain  $x$  or  $y$  which implies  $Y = \{y\}$  (because the vertices in  $Y$  are complete for  $G_2$ ). But, then  $x$  and  $y$  have comparable neighbourhood (Lemma 5).
- Let  $\omega([G_2]_G) \leq \omega - 2$ 
  - If  $|Y| \geq 2$  then any  $\omega$ -clique containing  $v \in G_2$ , contains  $y$  which is impossible (Proposition 2).
  - If  $Y = \{y\}$  then  $x$  and  $y$  have comparable neighbourhood (Lemma 5).

Lastly, pick up a vertex  $y$  in  $Y$  such that  $N(y) - S$  is maximal. We know that  $\forall y' \in Y$  such that  $yy' \notin E$ ,  $N_{G_1}(y') \subseteq N_{G_1}(y)$  (Lemma 22) and then  $\{y\} \cup N(y)$  forms a star-cutset. This completes the proof of the claim.  $\square$

*Case 2:  $q = 2$ .* If there exists an index  $i$  ( $1 \leq i \leq p$ ) such that  $A_i$  contains a vertex  $x$  complete for  $G_1 \cup G_2$  then  $\{x\} \cup G_1 \cup G_2 \cup V(B)$  forms a star-cutset, except if  $A_i = \{x\}$  and  $p = 2$  but this case cannot occur (Claim 25). Thus, we can assume that, for each  $i$  ( $1 \leq i \leq p$ ),  $A_i = A_i^1 \cup A_i^2$  (where  $A_i^1 \subseteq A_i$  is the subset of vertices in  $A_i$  which are incomplete for  $G_1$ ,  $A_i^2 = A_i \setminus A_i^1$ ).

**Claim 26.**  $\forall i$  ( $1 \leq i \leq p$ )  $A_i^1 = \emptyset$  or  $A_i^2 = \emptyset$ .

Let  $i$  be an index such that  $A_i^1 \neq \emptyset$  and  $A_i^2 \neq \emptyset$ . Let  $a \in A_i^1$  and  $b \in A_i^2$  such that  $ab \in E$  (they exist because  $A_i$  is a connected component of  $A$ ). Let  $s \in G_1 \setminus N_{G_1}(a)$  and  $r \in G_2 \setminus N_{G_2}(b)$ . Then for every  $u$  in  $A_j$  ( $j \neq i$ ) the subset  $\{r, a, b, s, u\}$  induces a  $P_5$  or a  $C_5$ .  $\square$

**Claim 27.**  $p = 2$ .

Assume that  $p \geq 3$ , so there exist two indices  $i$  and  $j$  ( $i \neq j$ ) such that  $A_i^1 \neq \emptyset$  and  $A_j^1 \neq \emptyset$  (or  $A_i^2 \neq \emptyset$  and  $A_j^2 \neq \emptyset$ ). Let  $x \in A_i^1 \cup A_j^1$  such that  $N(x) \cap G_1$  is maximal (say

$x \in A_i^1$ ), then

$$\forall y \in A_j^1, N_{G_1}(y) \subseteq N_{G_1}(x) \quad (\text{Lemma 22 with } v \in G_2).$$

But in this case,  $\{x\} \cup N_{G_1}(x) \cup G_2 \cup V(B)$  forms a star-cutset. So,

$$p = 2 \text{ and } A_1 = A_1^1, \quad A_2 = A_2^2. \quad \square$$

**Claim 28.**  $A_1$  and  $A_2$  induce complete subgraphs.

Assume that  $A_1$  (or  $A_2$ ) does not induce a complete subgraph. Let  $x \in A_1$  be a vertex such that  $x$  is not adjacent to all the other vertices in  $A_1$  and such that  $N(x) \cap G_1$  is maximal for inclusion. Let  $T \subset A_1$  be the subset of vertices in  $A_1$  which are not adjacent to  $x$ , then

$$\forall t \in T, \quad N_{G_1}(t) \subseteq N_{G_1}(x) \quad (\text{Lemma 22 with } v \in G_2).$$

But  $\{x\} \cup N_{G_1}(x) \cup G_2 \cup N_{A_1}(x) \cup V(B)$  forms a star-cutset. This completes the proof of the claim.  $\square$

Now we assume that  $|A_1| \geq |A_2|$ .

- (1)  $|A_2| \geq 2$ . Let  $v, w \in A_2$  and let  $K$  be an  $\omega$ -clique of  $G$  containing  $v$  but not  $w$  (Proposition 2).
  - If  $K \subseteq G_1 \cup A_2 \cup V(B)$  then one can find a clique of size  $\omega + 1$  by adding vertex  $w$  to  $K$  ( $w$  is adjacent to all vertices from  $G_1, A_2$  and  $V(B)$ ).
  - If  $K \subseteq G_2 \cup A_2 \cup V(B)$ , let  $K'$  be the clique formed by  $K \cap G_2$ ,  $K \cap V(B)$  and  $A_1$  (all the vertices in  $A_1$  are adjacent to all vertices from  $G_2$  and  $V(B)$ ), then  $K'$  is a clique of size strictly greater than  $\omega$  (since  $|A_1| > |A_2| - 1 \geq |K \cap A_2|$ ).
- (2)  $|A_2| = 1$ . We know that  $A_1$  is a clique and that all vertices in  $A_1$  are adjacent to all vertices from  $G_2$  and  $V(B)$ , so, if  $|A_1| \geq 2$  we have  $\omega([G_2]_G) \leq \omega - 2$ . Then let  $v \in G_2$  and  $x \in A_1$ , we cannot find an  $\omega$ -clique that contains  $v$  and not  $x$  which contradicts Proposition 2.

So  $|A_1| = |A_2| = 1$  (say  $A_1 = \{a_1\}$  and  $A_2 = \{a_2\}$ ). But in this case  $\{a_1, a_2\}$  forms an even-pair which contradicts Lemma 3.  $\square$

**Theorem 29** (Seinsche [14]). *A graph  $G$  is  $P_4$ -free if and only if for all induced subgraph  $H$  of  $G$  with more than one vertex,  $H$  or  $\bar{H}$  is disconnected.*

**Proof of Theorem 13.** Assume that  $[S]_G$  is  $P_4$ -free, since  $|S| \geq 2$  and  $[S]_G$  is connected (Theorem 12), we know that  $[S]_{\bar{G}}$  is disconnected (Theorem 29). There exists a partition of  $S$  in two nonempty subsets  $A$  and  $B$  such that

$$\forall a \in A, \quad \forall b \in B, \quad ab \in E.$$

Let  $S = A \cup B$  be such a partition (with  $A$  taken as small as possible).

- If  $|A| = 1$ , then  $S$  forms a star-cutset.



- If  $|A| \geq 2$ , then the subgraph induced by subset  $A$  is disconnected (else  $A$  is not minimal) which contradicts Lemma 10.  $\square$

**Proof of Theorem 19.** Since  $G$  is a minimal imperfect and  $(P_5, C_5)$ -free graph, we know that  $[N(v)]_G$  is connected (Corollary 14). Let  $w$  be a vertex of  $N(v)$  such that  $\mathcal{M}(w)$  is maximal for inclusion. Let  $Y = N(v) \setminus (\{w\} \cup N_{N(v)}(w))$ . We have  $Y \neq \emptyset$ , else any  $\omega$ -clique containing  $v$  also contains  $w$ , which is impossible. Now, we show that  $M = \mathcal{M}(w)$ . We know that  $\forall y \in Y \mathcal{M}(y) \subseteq \mathcal{M}(w)$  (Lemma 22), so if  $M \setminus \mathcal{M}(w) \neq \emptyset$ , the set  $\{w\} \cup N_{N(v)}(w) \cup \mathcal{M}(w)$  is a star disconnecting subgraphs induced by subsets  $M \setminus \mathcal{M}(w)$  and  $Y$ . We also have  $[Y]_G$  is connected else  $\{v\} \cup \{w\} \cup N_{N(v)}(w) \cup \mathcal{M}(w)$  is a star-cutset.  $\square$

Before proving Lemma 21, we recall some properties of bipartite connected graphs without induced  $P_5$ .

Let  $A$  be a connected bipartite graph on  $V(A) = V_1 \cup V_2$ . Let  $x \in V_1$ , for short we put  $N_2(x) = N_{V_2}(x)$ . The following three properties are equivalent:

- (1)  $\forall x, x' \in V_1, N_2(x) \subseteq N_2(x')$  or  $N_2(x') \subseteq N_2(x)$ .
- (2)  $A$  contains no induced  $P_5$ .
- (3)  $\forall y, y' \in V_2, N_1(y) \subseteq N_1(y')$  or  $N_1(y') \subseteq N_1(y)$ .

**Remark 30.** Property (1) shows that neighbourhoods, in  $V_2$ , of vertices in  $V_1$  are pairwise comparable which allow us to order them such that

$$N_2(x_i) \subseteq N_2(x_j) \quad \text{iff } i \geq j;$$

we do the same for vertices in  $V_2$ . One can remark that  $N_2(x_1) = V_2$  ( $A$  is connected).

We denote by  $\mathcal{U}_1$  (resp.  $\mathcal{U}_2$ ) the subset of vertices in  $V_1$  (resp.  $V_2$ ) which are universal for  $V_2$  (resp.  $V_1$ ) and  $\overline{\mathcal{U}}_1$  (resp.  $\overline{\mathcal{U}}_2$ ) the other vertices in  $V_1$  (resp.  $V_2$ ).

As before, let  $G$  be a graph and let  $S$  be a minimal cutset of  $G$ . We denote by  $G_1, G_2, \dots, G_q$  ( $q \geq 2$ ) the connected components of  $G \setminus S$  and we say that a vertex in  $S$  is

- of type U, if it is complete for  $G_1, G_2, \dots, G_q$ ;
- of type  $G_i$ , if it is incomplete for  $G_i$  ( $1 \leq i \leq q$ ).

**Lemma 31.** Let  $G$  be a minimal imperfect and  $(P_5, C_5)$ -free graph and let  $S$  be a minimal cutset of  $G$ . If  $xy$  is an edge of  $[S]_G$  such that  $x$  is of type  $G_i$  and  $y$  of type  $G_j$  ( $i \neq j$ ) then for every vertex  $z \in S$  we have  $xz \in E$  or  $yz \in E$ .

**Proof.** Assume there exists a vertex  $z$  in  $S$  such that  $xz \notin E$  and  $yz \notin E$ . Let  $u \in G_i \setminus N_{G_i}(x)$  and  $v \in G_j \setminus N_{G_j}(y)$  then the subset  $\{u, y, x, v, z\}$  induces a  $P_5$  or a  $C_5$ .  $\square$

**Proof of Lemma 21.** We assume that  $A$  (with  $V(A) = V_1 \cup V_2$ ) is a bipartite graph and that  $S$  is the complete join of  $A$  and  $B = S \setminus A$ . We know that  $A$  is connected (Lemma 10) and we claim that  $A$  is not a complete bipartite graph.

Indeed, if  $A$  is a complete bipartite graph (i.e.  $A$  is the complete join of two graphs  $A_1$  and  $A_2$  without edges, where  $V(A) = V(A_1) \cup V(A_2)$ ) then:

- If  $|V(A_1)| = 1$ , then  $S$  forms a star-cutset.
- If  $|V(A_1)| \geq 2$ , then  $S$  is the complete join of the two graphs  $A_1$  and  $A_2 \cup B$  (where  $G_1 \cup G_2 = (V(G_1) \cup V(G_2), E(G_1) \cup E(G_2))$ , for two graphs  $G_1$  and  $G_2$ ) but  $V(A_1)$  is a stable set which contradicts Lemma 10.

We order vertices in  $A$  like in Remark 30. We know that  $\mathcal{U}_1 \neq \emptyset$  (resp.  $\mathcal{U}_2 \neq \emptyset$ ) and since the bipartite graph is not complete we have  $\overline{\mathcal{U}_1} \neq \emptyset$  (resp.  $\overline{\mathcal{U}_2} \neq \emptyset$ ).

**Claim 32.** •  $|\mathcal{U}_1| \leq 2$  and  $|\mathcal{U}_2| \leq 2$  if  $q = 2$ .

- $|\mathcal{U}_1| = |\mathcal{U}_2| = 1$  if  $q \geq 3$ .

The vertices in  $\mathcal{U}_1$  cannot be of type U, otherwise they would have comparable neighbourhoods with vertices in  $\overline{\mathcal{U}_1} \neq \emptyset$  (Lemma 5).

If  $q \geq 3$  and  $|\mathcal{U}_1| \geq 2$ , let  $x_1 \in \mathcal{U}_1$  (resp.  $x_2 \in \mathcal{U}_1$ ) be of type  $G_i$  (resp.  $G_j$ ). Since  $q \geq 3$ , there exists  $k \neq i, j$  such that  $x_1$  and  $x_2$  are complete for  $G_k$ . Thus Lemma 22 with  $v \in G_k$  implies that  $i = j$ , but in this case  $x_1$  and  $x_2$  would have comparable neighbourhoods.

If  $q = 2$ , two distinct vertices (in  $\mathcal{U}_1$ ) would be of different types (i.e. incomplete for different components), else they would be comparable. So, in this case,  $|\mathcal{U}_1| \leq 2$ .  $\square$

**Claim 33.** If  $x \in \mathcal{U}_1$  (resp.  $\mathcal{U}_2$ ) is complete for  $G_i$ , then every vertex  $v$  in  $\overline{\mathcal{U}_2}$  (resp.  $\overline{\mathcal{U}_1}$ ) is complete for  $G_i$ .

Otherwise, assume that  $v \in \overline{\mathcal{U}_2}$  is not complete for  $G_i$ , then there exists  $x' \in \overline{\mathcal{U}_1}$  such that  $x'v \notin E$  and  $x'x \notin E$  which contradicts Lemma 31.  $\square$

**Claim 34.**  $|\overline{\mathcal{U}_1}| \leq 2$  and  $|\overline{\mathcal{U}_2}| \leq 2$

Let  $x \in \mathcal{U}_1$  be complete for  $G_i$ , then there exist at most two vertices in  $\overline{\mathcal{U}_2}$  which are complete for  $G_i$ . Indeed, assume that  $\{y_1, y_2, y_3\} \subseteq \overline{\mathcal{U}_2}$  are complete for  $G_i$ . We know that the clique number of the graph induced by  $G_i \cup S$  is  $\omega$  (any vertex of  $G_i$  belongs to  $\omega$   $\omega$ -cliques); moreover, the clique number of the graph induced by  $G_i \cup V(B)$  is  $\omega - 2$  (since  $xy_i \in E(A)$  ( $i = 1, 2, 3$ ) and both  $x$  and  $y_i$  ( $i = 1, 2, 3$ ) are adjacent to all vertices from  $G_i$  and  $B$ ). But this implies that  $y_1, y_2, y_3$  are pairwise co-critical which contradicts Lemma 8 ( $G$  is Berge). So,  $|\overline{\mathcal{U}_2}| \leq 2$  because Claim 33 implies that every vertex in  $\overline{\mathcal{U}_2}$  is complete for  $G_i$ .  $\square$

**Claim 35.** *If  $|\mathcal{U}_1| = 2$  (resp.  $|\mathcal{U}_2| = 2$ ) then  $|\overline{\mathcal{U}_2}| = 1$  (resp.  $|\overline{\mathcal{U}_1}| = 1$ ) and this vertex is of type U.*

Indeed, since  $|\mathcal{U}_1| = 2$  we have  $q = 2$  (Claim 32) and  $\mathcal{U}_1$  contains a vertex complete for  $G_1$  and another complete for  $G_2$ . This implies that all vertices in  $\overline{\mathcal{U}_2}$  are of type U (Claim 33). But if  $|\overline{\mathcal{U}_2}| = 2$ , those two vertices have comparable neighbourhoods. This completes the proof of the claim.  $\square$

Now, we can finish the proof. First, we assume that  $\overline{\mathcal{U}_2}$  contains a vertex of type U (say  $y$ ). A maximal stable set containing  $y$  will be included in the bipartite subgraph induced by  $V(A)$  (because  $y$  is complete for  $G_1, G_2, \dots, G_q$  and  $B$ ). So this stable set will be of size at most  $|\overline{\mathcal{U}_1}| + |\overline{\mathcal{U}_2}|$  or  $|\mathcal{U}_2| + |\overline{\mathcal{U}_2}|$  (i.e. of size at most 4). If this stable set is of size 3 or less, we can conclude using Tucker's Theorem [15]. So, it remains two cases to study:

*Case 3  $|\overline{\mathcal{U}_1}| = |\overline{\mathcal{U}_2}| = 2$ :* We can suppose that there exist no edges (in  $G$ ) between vertices in  $\overline{\mathcal{U}_1}$  and vertices in  $\overline{\mathcal{U}_2}$  (else the stable set could not be of size 4). But the two vertices in  $\overline{\mathcal{U}_2}$  have comparable neighbourhoods, a contradiction.

*Case 4  $|\mathcal{U}_2| = |\overline{\mathcal{U}_2}| = 2$ :* There exist no edges between the vertex in  $\overline{\mathcal{U}_1}$  (Claim 35) and those in  $\overline{\mathcal{U}_2}$  (because vertices in  $\overline{\mathcal{U}_2}$  are not adjacent to all  $\mathcal{U}_1 \cup \overline{\mathcal{U}_1}$ ). In this case, the two vertices in  $\overline{\mathcal{U}_2}$  have comparable neighbourhoods, a contradiction.

So neither  $\overline{\mathcal{U}_1}$  nor  $\overline{\mathcal{U}_2}$  contain vertices of type U. This implies that (Claim 35)

$$|\mathcal{U}_1| = |\mathcal{U}_2| = 1.$$

We note  $\mathcal{U}_1 = \{x_1\}$ ,  $\mathcal{U}_2 = \{y_1\}$ ,  $\overline{\mathcal{U}_1} = \{x_2\}$  or  $\{x_2, x_3\}$ ,  $\overline{\mathcal{U}_2} = \{y_2\}$  or  $\{y_2, y_3\}$ ,  $V_1 = \mathcal{U}_1 \cup \overline{\mathcal{U}_1}$  and  $V_2 = \mathcal{U}_2 \cup \overline{\mathcal{U}_2}$ .

Suppose that  $x_1$  is of type  $G_i$ , then the vertices in  $\overline{\mathcal{U}_2}$  are of type  $G_i$  (Claim 33 and there exist no vertices of type U). If  $y_1$  is of type  $G_j$  ( $j \neq i$ ) then the vertices in  $\overline{\mathcal{U}_1}$  are of type  $G_j$ , moreover, if  $|\overline{\mathcal{U}_1}| = 2$  we have  $N_{G_i}(x_2) \subseteq N_{G_i}(x_3)$  (or  $N_{G_j}(x_3) \subseteq N_{G_j}(x_2)$ ) and  $N_{G_j}(x_{|V_1|}) \neq G_j$  (because there is no vertices of type U). But, then  $\{x_1\} \cup V_2 \cup V(B) \cup N_{G_j}(x_{|V_1|}) \cup \bigcup_{k \neq i, j} G_k$  forms a star disconnecting subgraphs induced by the sets  $G_j \setminus N_{G_j}(x_{|V_1|})$  and  $\overline{\mathcal{U}_1}$ . So,  $y_1$  and the vertices in  $\overline{\mathcal{U}_1}$  are of type  $G_i$  and all the vertices  $x_k$  ( $1 \leq k \leq |V_1|$ ) are complete for  $G_j$  ( $j \neq i$ ). Thus Lemma 22 with  $v \in G_j$  and the fact that no two vertices can have comparable neighbourhoods imply:  $N_{G_i}(x_k) \subset N_{G_i}(x_l)$  for  $k < l$ .

The same holds for the vertices  $y_k$ . We have  $x_{|V_1|}y_{|V_2|} \notin E$  and,  $x_{|V_1|}$  and  $y_{|V_2|}$  are of type  $G_i$ . So Lemma 22 (with  $v \in G_k$ ,  $k \neq i$ ) implies  $N_{G_i}(x_{|V_1|}) \subseteq N_{G_i}(y_{|V_2|})$  or  $N_{G_i}(y_{|V_2|}) \subseteq N_{G_i}(x_{|V_1|})$ , say  $N_{G_i}(y_{|V_2|}) \subseteq N_{G_i}(x_{|V_1|})$ . Since the vertex  $x_{|V_1|}$  is not complete for  $G_1, \dots, G_q$ , the subset  $\{x_{|V_1|}\} \cup N_{G_i}(x_{|V_1|}) \cup \bigcup_{j \neq i} G_j \cup V(B)$  is a star-cutset of  $G$ . This completes the proof of the Lemma.  $\square$

**Theorem 36.** *Let  $G$  be a  $P_5$ -free minimal imperfect Berge graph and let  $S$  be a minimal cutset of  $G$ . Then  $S$  cannot induce a graph in  $\mathcal{B}^*$ .*

**Proof.** Assume that  $[S]_G \in \mathcal{B}^*$ . We know that  $[S]_G$  is connected, then  $[S]_G$  is either a bipartite graph (which contradicts Lemma 21 with  $A = [S]_G$  and  $V(B) = \emptyset$ ) or the complete join of two graphs  $G_1$  and  $G_2$ . In the latter case, assume that  $[S]_G$  is the complete join of  $G_1$  and  $G_2$  such that  $G_1$  was not decomposable by the complete join operation. Then, either  $G_1$  is a bipartite graph (contradicts Lemma 21) or  $G_1$  is a disconnected subgraph of  $[S]_G$  (contradicts Lemma 10).  $\square$

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